

1a. From equation 2, the expected payoff of the game is

$$pAq = \begin{bmatrix} \frac{1}{2} & 0 & \frac{1}{2} \end{bmatrix} \begin{bmatrix} -4 & 6 & -4 & 1 \\ 5 & -7 & 3 & 8 \\ -8 & 0 & 6 & -2 \end{bmatrix} \begin{bmatrix} \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \end{bmatrix} = -5/8$$

b. If player R uses strategy  $[p_1 \ p_2 \ p_3]$  against player C's strategy  $[\frac{1}{4} \ \frac{1}{4} \ \frac{1}{4} \ \frac{1}{4}]^t$ , his payoff will be  $pAq = (-\frac{1}{4})p_1 + (\frac{9}{4})p_2 - p_3$ . Hence

$p_1, p_2, p_3$  are non-negative numbers and add up to 1, this is a weighted average of numbers  $-\frac{1}{4}, \frac{9}{4}$ , and  $-1$ . Clearly, this is largest if  $p_1 = p_3 = 0$  and  $p_2 = 1$ . so  $p = [0 \ 1 \ 0]$

c. As in (b), if player C uses  $[q_1 \ q_2 \ q_3 \ q_4]^t$  against  $[\frac{1}{2} \ 0 \ \frac{1}{2}]$ ,

$pAq = -6q_1 + 3q_2 + q_3 - \frac{1}{2}q_4$ . This is minimized by setting  $q_1 = 1$  and  $q_2 = q_3 = q_4 = 0$ . so,  $q = [1 \ 0 \ 0 \ 0]^t$

3a. Calling the matrix  $A$ , we see  $a_{22}$  is a saddle point, so the optimal strategies are pure;  $p = [0 \ 1]$ ,  $q = [0 \ 1]^t$ ; the value of the game is  $a_{22} = 3$ .

b. As in (a),  $a_{21}$  is a saddle point, so optimal strategies are  $p = [0 \ 1 \ 0]$ ,  $q = [1 \ 0]^t$ ; the value of the game is  $a_{21} = 2$ .

c. Here,  $a_{32}$  is a saddle point, so optimal strategies are  $p = [0 \ 0 \ 1]$ ,  $q = [0 \ 1 \ 0]^t$  and  $v = a_{32} = 2$ .

d. Here,  $a_{21}$  is a saddle point, so  $p = [0 \ 1 \ 0 \ 0]$ ,  $q = [1 \ 0 \ 0]^t$  and  $v = a_{21} = -2$ .

4a. Calling the matrix  $A$ , the formula from Theorem 2 on P.605 yields

$$p = \begin{bmatrix} 5/8 & 3/8 \end{bmatrix}, \quad q = \begin{bmatrix} 1/8 & 7/8 \end{bmatrix}^t, \quad v = 27/8. \quad A \text{ has no saddle points.}$$