

# MATHEMATICA ROUTINES FOR THE DIRAC OPERATOR FOR GRAPHS

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ABSTRACT. Mathematica routines for the Dirac operator. This allows to compute the cohomology for any graph. [1].

The first routine finds all the cliques in a graph.

```

Cli[s_, k_] := Module[{m, n, c, u, V, W, l = {}},
  V = VertexList[s]; n = Length[V];
  W = EdgeList[s]; m = Length[W];
  W = Table[{W[[j, 1]], W[[j, 2]]}, {j, Length[W]}];
  c = Subsets[V, {k, k}];
  If[k == 1, l = V, If[k == 2, l = W,
  Do[ss = Subgraph[s, c[[j]]];
    If[Length[EdgeList[ss]] == Binomial[k, 2],
    l = Append[l, VertexList[ss]],
    {j, Length[c]}]]; l];

```

Now we compute the matrix  $D$ .

```

Dirac[s_] := Module[{q, d, l, b, u, m, n, p, v, DD},
  q = VertexList[s]; n = Length[q];
  d = Table[{0}, {p, n - 1}];
  l = Table[{}], {p, n}; b = Table[0], {p, n}; m = n;
  Do[If[m == n,
    l[[p]] = Cli[s, p]; b[[p]] = Length[l[[p]]];
    If[b[[p]] == 0, m = p - 2], {p, n}];
  v = Sum[b[[p]], {p, n}];
  u = Table[Sum[b[[p]], {p, 1, k}], {k, Min[n, m + 1]}];
  u = Prepend[u, 0]; DD = Table[0, {v}, {v}];
  If[m > 0, d[[1]] = Table[0, {j, b[[2]]}, {i, b[[1]]}];
  Do[d[[1, j, 1[[2, j, 1]]]] = -1, {j, b[[2]]};
  Do[d[[1, j, 1[[2, j, 2]]]] = 1, {j, b[[2]]}];
  Do[If[m >= p,
    d[[p]] = Table[0, {j, b[[p + 1]]}, {i, b[[p]]}];
    Do[a = 1[[p + 1, i]];
    Do[k = Position[l[[p]], Delete[a, j]][[1, 1]];
    d[[p, i, k]] = (-1)^j, {j, p + 1},
    {i, b[[p + 1]]}], {p, 2, n - 1}];

```

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```

Do[
  If [m>=p, Do[DD[[u[[p+1]]+j, u[[p]]+i]] = d[[p, j, i]],
    {i, b[[p]]}, {j, b[[p+1]]}], {p, 1, n-1}];
{DD+Transpose[DD], u}];

```

This leads to the Laplace Beltrami operators which allow to compute the Betti numbers, and the Cohomology basis.

```

LaplaceBeltrami[s_]:=Module[{DD, LL, br},
  {DD, br}=Dirac[s]; LL=DD.DD;
  Table[Table[LL[[br[[k]]+i, br[[k]]+j]],
    {i, br[[k+1]]-br[[k]]}, {j, br[[k+1]]-br[[k]]},
    {k, Length[br]-1}];
Nulley[A_]:=Length[A]-MatrixRank[A];
Betti[s_]:=Map[Nulley, LaplaceBeltrami[s]];

```

Example:

```

(* This is an attempt to expand the Dirac code and comment *)
Cli[s_, k_]:=Module[{m, n, c, u, V, W, l={}},
  V=VertexList[s]; n=Length[V];
  W=EdgeList[s]; m=Length[W];
  W=Table[{W[[j, 1]], W[[j, 2]]}, {j, Length[W]}];
  c=Subsets[V, {k, k}];
  If[k==1, l=V,
  If[k==2, l=W,
  Do[
    ss = Subgraph[s, c[[j]]];
    If[Length[EdgeList[ss]]==Binomial[k, 2],
    l = Append[l, VertexList[ss]],
    {j, Length[c]}];
  ]];
  l];
  (* will return list of subgraphs *)
  (* Vertices and length *)
  (* Edges and length *)
  (* write edges as pairs *)
  (* all the subsets of length k *)
  (* in case k=1, take vertices *)
  (* in case k=2, take edges *)
  (* start to sum over all subsets *)
  (* build subgraph *)
  (* is it complete? *)
  (* if yes, add it *)
  (* end sum over subsets *)
  (* end if statements *)
  (* return collection of subsets *)

Dirac[s_]:=Module[{q, d, l, b, bb, m, n, p, v, DD},
  q=VertexList[s]; n=Length[q];
  d=Table[{0}, {p, n-1}];
  l=Table[{0}, {p, n}];
  b=Table[0, {p, n}];
  m=n;
  Do[If[m==n,
    l[[p]]=Cli[s, p];
    b[[p]]=Length[l[[p]]];
    If[b[[p]]==0, m=p-2], {p, n}];
  v=Sum[b[[p]], {p, n}];
  bb=Table[Sum[b[[p]], {p, 1, k}], {k, Min[n, m+1]}];
  bb=Prepend[bb, 0];
  DD=Table[0, {v}, {v}];
  If[m>0,
    d[[1]] = Table[0, {j, b[[2]]}, {i, b[[1]]}];
    Do[d[[1, j, l[[2, j, 1]]]] = -1, {j, b[[2]]};
    Do[d[[1, j, l[[2, j, 2]]]] = 1, {j, b[[2]]};
  ];
  Do[
    If [m>=p,
    d[[p]] = Table[0, {j, b[[p+1]]}, {i, b[[p]]}];
    Do[
      a=1[[p+1, i]];
      Do[
        k=Position[l[[p]], Delete[a, j]][[1, 1]];
        d[[p, i, k]] = (-1)^j,
        {j, p+1};
      {i, b[[p+1]]};
    ],
    {p, 2, n-1};
  Do[
    If [m>=p,
    Do[
      DD[[bb[[p+1]]+j, bb[[p]]+i]] = d[[p, j, i]],
      {i, b[[p]]},
      {j, b[[p+1]]}],
    {p, 1, n-1}];
  (* Find simplices and lengths *)
  (* Vertex list and its length *)
  (* Container for d.k submatrices *)
  (* Container for cliques *)
  (* Container for cliques numbers *)
  (* upper bound of clique dim *)
  (* as long as there are cliques *)
  (* build list of all cliques *)
  (* find its lengths *)
  (* abort if there are none *)
  (* this is the size of the matrix *)
  (* these are the subdivisions *)
  (* have also the left subdivision *)
  (* empty shell for the matrix d *)
  (* if m>0, we have some edges *)
  (* fill in empty shell *)
  (* place the -1 entries in *)
  (* place the 1 entries in *)
  (* we have finished with edges *)
  (* now do the same for triangles *)
  (* only if not reached max *)
  (* dk matrices *)
  (* core of the procedure *)
  (* the clique we work on *)
  (* start run over all subsimplex *)
  (* find label of the subsimplex *)
  (* orientation of subsimplex *)
  (* run over all subsimplices *)
  (* run over simplices of dim p *)
  (* finish if condition *)
  (* sum over dimensions *)
  (* now fill in the d.k matrices *)
  (* only if dimension small enough *)
  (* start summing *)
  (* glue in the matrix *)
  (* sum over rows *)
  (* sum over columns *)
  (* sum over dimensions *)

```

```
DD+Transpose[DD];
```

```
(* finally build d+d^*
```

```
*)
```

## REFERENCES

- [1] O. Knill. The McKean-Singer Formula in Graph Theory.  
<http://arxiv.org/abs/1301.1408>, 2012.  
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